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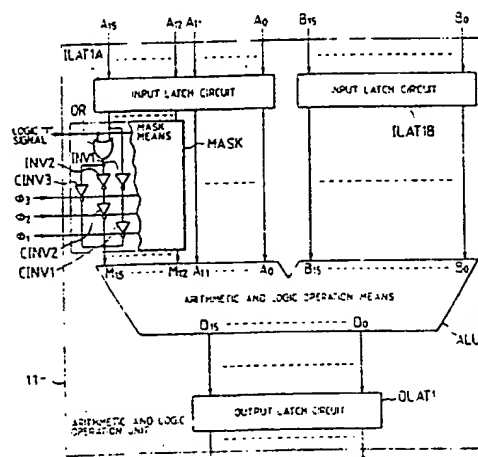
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⑤4 Data processor.

57) A processor capable of processing data of a structure having a first information field and a second information field includes conservation means for conserving the information of the predetermined second information field of the input data. Thus, in executing a logical operation or a shift operation, the step of separating the first data fields and second data fields from the data to-be-processed and the step of afficing the predetermined second field to the result of the operation are dispensed with, thereby to achieve a high-speed operation for the data of the structure having the first information field and the second information field.

**EP 0 307 166 A2**

FIG. 1



## DATA PROCESSOR

The present invention relates to a data processor which includes operation means, such as an arithmetic and logic operation unit or a shift operation unit, capable of processing data of a structure having a first information field and a second information field. By way of example, it relates to a technique which is effective when applied to the operations of abstract data each having a data field and a tag field for expressing the type or attribute of data contained in the data field.

Abstract data in which the attribute or data type of data such as the integer type or the floating point type is expressed by a tag field is, for example, one of a data structure having a data field and the tag field within one word. An arithmetical or logical operation or a shift operation for such abstract data needs to be executed, in effect, for the information of the data fields, and the operated result is set as the information of the data field pairing with the tag field expressing the attribute thereof.

Heretofore, an arithmetic and logic operation unit or a barrel shifter has been constructed so as to execute an arithmetical or logical operation or a shift operation for the respective bits of data received as inputs. Accordingly, in a case where only partial fields contained in data to-be-processed are subjected to a substantial operation because of the property of the data structure thereof as in the abstract data, there must be adopted a procedure wherein data fields and tag fields are separated from the data to-be-processed, the information items of the separated data fields are subjected to a predetermined operation, and a predetermined tag field is affixed to the operated result. Incidentally, an example of a literature in which a data processor is described is "LSI Handbook," p. 536, published by the Ohm-Sha, Ltd. on November 30, 1984.

However, in the case of executing the operation of the abstract data in accordance with the above procedure, there has been the problem that the speed of the operation lowers drastically due to the necessity for the step of separating the data fields and the tag fields from the data to-be-processed, the step of subjecting the information items of the separated data fields to the predetermined operation, and the step of affixing the predetermined tag field to the operated result. Such a problem forms an obstacle to raising the speed of a process which is based on an AI (artificial intelligence) language and which proceeds while deciding the types of data during the run of a program.

Preferred embodiments of the present inven-

tion provide a data processor which can execute operations at high speed for data of a structure having a first information field and a second information field.

The above and other desirable features of the present invention will become apparent from the description of this specification and the accompanying drawings.

Typical aspects of performance of the present invention are briefly summarized as follows:

An arithmetic and logic operation unit which can accept data items on two input sides and then process them, each of the data items being of a structure having a first information field and a second information field, is provided on one of the input sides with mask means capable of masking the second information field contained in the input data, into a predetermined logical value in accordance with the sort of an arithmetical or logical operation. By way of example, in a case where a logical sum, exclusive "or" or addition operation or the like is executed by the arithmetic and logic operation unit, the second information field of the input data on one side is masked into logic "0" by the mask means, and in a case where a logical product or subtraction operation or the like is executed, the second information field of the input data on one side is masked into logic "1" by the mask means. In this way, the arithmetical or logical operation for one pair of input data is executed, and the operated result contains the result of the arithmetical or logical operation between the first information fields and also conserves the information of the second information field of the input data on the other side as left intact.

In addition, a shift operation unit, such as barrel shifter, which can accept data items on two input sides and then subject them to a shift operation, each of the data items being of a structure having a first information field and a second information field, is constructed of rearrange means capable of rearranging the information fields so that information items of the first information fields respectively contained in one pair of input data to-be-processed may be put into a serial form, shift operation means for executing the shift operation for data delivered as an output from the rearrange means, and replace means capable of replacing a part of data delivered as an output from the shift operation means, with information of the second information field contained in the input data on one side. According to this shift operation unit, the information items of the first information fields respectively contained in one pair of input data to-be-processed are rearranged so as to become interlacingly serial

and then subjected to the shift operation, and the operated result is partly replaced with the information of the predetermined second information field, which is conserved.

A data processor including the arithmetic and logic operation unit or the shift operation unit as described above makes it unnecessary, in executing the logical operation or the shift operation, to especially perform the step of separating the first data fields and second data fields from the data to-be-processed, as well as the step of affixing the predetermined second field to the result of the operation, thereby to achieve a high-speed operation for the data of the structure having the first information field and the second information field.

### BRIEF DESCRIPTION OF THE DRAWINGS

Fig. 1 is a block diagram of an embodiment of an arithmetic and logic operation unit according to the present invention;

Fig. 2 is a diagram for explaining the operation of the embodiment in Fig. 1 in the case of masking a tag field into logic "1" in the arithmetical or logical operation for data with tags;

Fig. 3 is a diagram for explaining the operation of the embodiment in Fig. 1 in the case of masking a tag field into logic "0" in the arithmetical or logical operation for data with tags;

Fig. 4 is a diagram for explaining the arithmetical or logical operation of the embodiment in Fig. 1 for data with no tag;

Fig. 5 is a block diagram of an embodiment of a shift operation unit according to the present invention;

Fig. 6 is a logic diagram of an example of rearrange means included in the embodiment in Fig. 5;

Fig. 7 is a block diagram of an example of replace means included in the embodiment in Fig. 5;

Fig. 8 is a diagram for explaining the shift operation of the embodiment in Fig. 5 for data with tags;

Fig. 9 is a diagram for explaining the shift operation of the embodiment in Fig. 5 for data with no tag;

Fig. 10 is a block diagram of an embodiment of an execution unit which comprises the arithmetic and logic operation unit shown in Fig. 1 and the shift operation unit shown in Fig. 5; and

Fig. 11 is a block diagram of an embodiment of a processor which includes the execution unit shown in Fig. 10.

Shown in Fig. 11 is a coprocessor which is an embodiment of a data processor according to the

present invention. Although not especially restricted, the coprocessor 1 illustrated in the figure is formed on a single semiconductor substrate such as silicon substrate by known techniques for manufacturing semiconductor integrated circuits.

The coprocessor 1 serves to assist the processing ability of an external main processor, not shown, coupled thereto through a bus interface circuit 2 or to lighten the burden of processes on the main processor. It executes processes in accordance with the instructions of the main processor. In particular, this coprocessor 1 includes an execution unit 3 for processing under microprogram controls both abstract data (hereinbelow, also simply termed "data with tags") of a structure having a data field as a first information field and a tag field as a second information field for expressing the attribute or type of data contained in the data field, and data (hereinbelow, also simply termed "data with no tag") of a structure made up of only a data field. Here, although not especially restricted, the tag field is used as a field expressing the attribute or type of data, such as the integer type or the floating point type.

The coprocessor 1 of this embodiment comprises a microinstruction ROM (read-only memory) 4 which stores therein microprograms descriptive of various processing steps by the execution unit 3, etc. The microinstruction ROM 4 is accessed by a controller 5, whereby microinstructions constituting the microprogram are successively read out.

The controller 5 fetches a command given by the unshown main processor through the bus interface circuit 2 as well as an internal bus 6, and it accesses the microinstruction ROM 4 on the basis of address signals obtained by decoding a command code contained in the command or on the basis of address information directly contained in the command. Thus, the first one of a series of microinstructions for executing a process designated by the command is read out. Each of the second one et seq. of the series of microinstructions for executing the process instructed by the command is designated in such a way that the information of a next-address field in the microinstruction read out immediately before is supplied to the controller 5. The microinstruction read out of the microinstruction ROM 4 in this way is supplied to a microinstruction decoder 7. This microinstruction decoder 7 decodes the given microinstruction, and generates control signals 8 for the execution unit 3, etc. By the way, when the branch of a microinstruction flow has become necessary during the execution of the process or the like according to the microinstructions by the execution unit 3, this execution unit 3 gives the controller 5 an instruction therefor.

The execution unit 3 is coupled to the internal

bus 6, and is also coupled to a pair of RAMs (random access memories) 9 and 10. The RAMs 9 and 10 store externally-supplied data necessary for operations beforehand, and they are utilized as temporary registers in cases of processes. Accesses to the RAMs 9 and 10 are controlled on the basis of the control signals 8 which are output from the microinstruction decoder 7. By way of example, when the predetermined areas of the RAMs 9 and 10 are utilized as the temporary registers in case of a floating-point operation, the RAMs 9 and 10 are subjected to such a read-modify-write operation that source data is read out and that operated result data for the read source data is brought back into the RAMs 9 and 10 as destination data.

Fig. 10 shows an example of the execution unit 3.

Although not especially restricted, this execution unit 3 includes internal buses BUS1, BUS2 and BUS3 each of which is composed of 16 bits. Coupled to the internal buses BUS1, BUS2 and BUS3 are an arithmetic and logic operation unit 11, a shift operation unit 12, a shift counter SCUNT and a temporary register 13 which are typically illustrated. Incidentally, the temporary register 13 can be assigned to the predetermined areas of the RAMs 9 and 10.

Fig. 1 exemplifies the arithmetic and logic operation unit 11.

This arithmetic and logic operation unit 11 has input latch circuits ILAT1A and ILAT1B which are respectively supplied with 16-bit data  $A_0 - A_{15}$  on one side and 16-bit data  $B_0 - B_{15}$  on the other side to-be-processed. The data items  $A_0 - A_{15}$  and  $B_0 - B_{15}$  to be processed are supplied from the internal buses BUS1 and BUS2 separately from each other. In this regard, in a case where the data  $A_0 - A_{15}$  is supplied from the internal bus BUS1 to the input latch circuit ILAT1A, the data  $B_0 - B_{15}$  to be given to the input latch circuit ILAT1B is supplied from the other internal bus BUS2. On the other hand, in a case where the data  $A_0 - A_{15}$  is supplied from the internal bus BUS2 to the input latch circuit ILAT1A, the data  $B_0 - B_{15}$  to be given to the input latch circuit ILAT1B is supplied from the internal bus BUS1. Which of the internal buses the required data is to be accepted from, is controlled by a selection gate, not shown, which is controlled on the basis of the designation of the microinstruction.

Although no special restriction is meant, when one pair of abstract data are to be processed in the arithmetic and logic operation unit 11 of this embodiment, the information of the tag field to be conserved in a processed result is set to be the information of the tag field which is latched in one input latch circuit ILAT1B. In other words, data which contains the information of the tag field to be conserved in the processed result is supplied to

the input latch circuit ILAT1B. In this embodiment, the tag field which expresses the attribute or data type of data such as the integer type or the floating point type is composed of copper 4 bits within 1 word (16 bits) though not especially restricted.

In Fig. 1, symbol MASK denotes mask means capable of masking the upper 5 bits  $A_{15} - A_{12}$  delivered as outputs from the input latch circuit ILAT1A, into a predetermined logical value in accordance with the sort of a logical operation and the sort of the input data. More specifically, as regards the data made up of only the data field, the input bits are directly output without being masked. As regards the data containing the tag field, the upper 4 bits constituting the tag field are masked into logic "0" for a logical sum, exclusive "or" or addition operation, and they are masked into logic "1" for a logical product or subtraction operation.

The mask means MASK has the same circuit arrangements as one another for the respective bits  $A_{15} - A_{12}$ . For example, regarding the bit  $A_{15}$ , the circuit arrangement includes a path along which the bit  $A_{15}$  is output with its logical value left intact by an inverter INV1 and a clocked inverter CINV1 that are connected in series, a path along which the bit  $A_{15}$  is altered into logic "1" and then output by an inverter INV2 and a clocked inverter CINV2 that are connected in series to the output terminal of a 2-input "or" gate OR normally supplied with a signal of logic "1" and the bit  $A_{15}$ , and a path along which the bit  $A_{15}$  is altered into logic "0" and then output by a clocked inverter CINV3 that is connected in series to the output terminal of the "or" gate OR.

The selection terminals of the clocked inverters CINV1 - CINV3 are respectively supplied with control signals  $\phi_1$ ,  $\phi_2$  and  $\phi_3$ , whereby any of the paths is selected for the input bits  $A_{15} - A_{12}$ .

Although not especially restricted, the control signals  $\phi_1$ ,  $\phi_2$  and  $\phi_3$  are output from the microinstruction decoder 7 which decodes the microinstructions read out of microinstruction ROM 4 in succession. The series of microinstructions for each of various arithmetical and logical operations as stored in the microinstruction ROM 4 contain information for specifying the sort of the corresponding operation, and information for designating whether or not data to-be-processed is abstract data. Which of the control signals  $\phi_1$ ,  $\phi_2$  and  $\phi_3$  is to be brought to an operating select level is determined in accordance with the above information items.

In Fig. 1, symbol ALU denotes an arithmetic and logic operation means capable of executing the respective operations of logical sum, exclusive "or", addition, logical product and subtraction though not especially restricted. The input terminals of the arithmetic and logic operation means

ALU on one side are supplied with 4 bits  $M_{15} - M_{12}$  delivered as outputs from the mask means MASK and the lower 12 bits  $A_{11} - A_0$  delivered as outputs from the input latch circuit ILAT1A, while the input terminals thereof on the other side are supplied with the 16 bits  $B_{15} - B_0$  delivered as outputs from the input latch circuit ILAT1B. The operations of logical sum, exclusive "or", etc. in the arithmetic and logic operation means ALU are controlled by the output control signals 8 of the microinstruction decoder 7. Operated result data items  $D_{15} - D_0$  are selectively given to the internal buses BUS1, BUS2 and BUS3 through an output latch circuit OLAT1.

Here, the tag field to be conserved in the arithmetic and logic operation unit 11 of this embodiment is set at the bits  $B_{15} - B_{12}$  as stated above, and the bits  $A_{15} - A_{12}$  are masked into the logical value "1" or "0" in order to conserve the tag field in the case of the arithmetical or logical operation. On this occasion, in the case of the arithmetical or logical operation for the abstract data, no carry from the lower bits is transmitted to the operated result bits  $D_{15} - D_{12}$  which are the contents of the tag field to be conserved, and the bits  $B_{15} - B_{12}$  are conserved as the upper 4 bits  $D_{15} - D_{12}$  of the operated result without any change. By way of example, a gate capable of selectively blocking a path for transmitting a carry, which might occur on the basis of the operation between the input bits  $A_{11}$  and  $B_{11}$ , to the upper bits is provided though not shown. Such a gate can be constructed of an "and" gate, one input terminal of which is supplied with the control signal of logic "0" from the microinstruction decoder 7 in the case of the arithmetical or logical operation of the abstract data.

Next, the operation of the arithmetic and logic operation unit 11 will be described.

First, the operation in the case of subjecting the input data  $A_0 - A_{15}$  and  $B_0 - B_{15}$  as the abstract data to the logical product (AND) or subtraction (SUB) process will be described with reference to Fig. 2.

The input data items  $A_0 - A_{15}$  and  $B_0 - B_{15}$  have their upper 4 bits  $A_{15} - A_{12}$  and  $B_{15} - B_{12}$  set as tag fields TFa and TFb, respectively, and have their lower 12 bits  $A_0 - A_{11}$  and  $B_0 - B_{11}$  set as data fields DFa and DFb, respectively.

When it is instructed to execute the logical product or subtraction process for the data with the tags, microinstructions responsively read out of the microinstruction ROM 4 in succession are decoded by the microinstruction decoder 7, and the mask means MASK and the arithmetic and logic operation means ALU are controlled. More specifically, the control signal  $\phi_2$  is brought to the select level, the respective bits  $A_{15} - A_{12}$  of the tag field TFa in

the input data  $A_0 - A_{15}$  are masked into logic "1", and data containing the bits masked into such a logical value and the input data  $B_0 - B_{15}$  on the other side are subjected to the logical product operation or subtraction operation by the arithmetic and logic operation means ALU. In this way, the information items  $D_{15} - D_{12}$  of a tag field TFe in the operated result data  $D_0 - D_{15}$  are brought to the logical product or subtraction results between "1, 1, 1, 1" and " $B_{15}, B_{14}, B_{13}, B_{12}$ ", and the information items  $B_{15} - B_{12}$  of the tag field TFb in the input data  $B_0 - B_{15}$  are conserved. The information items  $D_{11} - D_0$  of a data field DFe in the operated result data  $D_0 - D_{15}$  are brought to the logical product or subtraction results between  $A_{11} - A_0$  and  $B_{11} - B_0$ .

Next, the operation in the case of subjecting the input data  $A_0 - A_{15}$  and  $B_0 - B_{15}$  as the abstract data to the logical sum (OR), exclusive "or" (EOR) or addition (ADD) process will be described with reference to Fig. 3.

In this case, microinstructions successively read out of the microinstruction ROM 4 in conformity with the sort of such an operation are decoded by the microinstruction decoder 7, whereby the control signal  $\phi_3$  is brought to the select level, the respective bits  $A_{15} - A_{12}$  of the tag field TFa in the input data  $A_0 - A_{15}$  are masked into logic "0", and data containing the bits masked into such a logical value and the input data  $B_0 - B_{15}$  on the other side are subjected to the logical sum operation, exclusive "or" operation or addition operation by the arithmetic and logic operation means ALU. In this way, the information items  $D_{15} - D_{12}$  of the tag field TFe in the operated result data  $D_0 - D_{15}$  are brought to the logical sum, exclusive "or" or addition results between "0, 0, 0, 0" and " $B_{15}, B_{14}, B_{13}, B_{12}$ ", and the information items  $B_{15} - B_{12}$  of the tag field TFb in the input data  $B_0 - B_{15}$  are conserved. The information items  $D_{11} - D_0$  of the data field DFe in the operated result data  $D_0 - D_{15}$  are brought to the logical sum, exclusive "or" or addition results between  $A_{11} - A_0$  and  $B_{11} - B_0$ .

Accordingly, in executing the arithmetical and logical operations for the data with the tags as shown in Figs. 2 and 3, the step of separating the data fields and tag fields from the data to-be-processed and the step of affixing a predetermined tag field to an operated result for the information items of the separated data fields are dispensed with, so that high-speed operations for the data with the tags are achieved.

In executing an arithmetical or logical operation for input data  $A_0 - A_{15}$  and  $B_0 - B_{15}$  made up of only data fields as shown in Fig. 4, microinstructions successively read out of the microinstruction ROM 4 in conformity with the sort of such an operation are decoded by the microinstruction de-

coder 7, whereby the control signal  $\phi_1$  is brought to the select level. Thus, the input data items  $A_0 - A_{15}$  and  $B_0 - B_{15}$  are supplied to the arithmetic and logic operation means ALU without any change and are used for the predetermined logical operation.

Fig. 5 exemplifies the shift operation unit 12 such as barrel shifter.

Likewise to the arithmetic and logic operation unit 11, the shift operation unit 12 shown in Fig. 5 executes shift operations for both data with tags, each having the tag field assigned to the upper 4 bits in 1 word, and data with no tag.

This shift operation unit 12 has input latch circuits ILAT2A and ILAT2B which are respectively supplied with 16-bit data  $A_0 - A_{15}$  on one side and 16-bit data  $B_0 - B_{15}$  on the other side to-be-processed, through the internal buses BUS1 and BUS2. A method of supplying the data from the internal buses BUS1 and BUS2 to the input latch circuits ILAT2A and ILAT2B is the same as in the case of the arithmetic and logic operation unit 11. In a case where the input data items are the data with the tags, the upper 4 bits  $A_{15} - A_{12}$  and  $B_{15} - B_{12}$  constitute tag fields TFa and TFb, respectively, and the lower 12 bits  $A_0 - A_{11}$  and  $B_0 - B_{11}$  constitute data fields DFa and DFb, respectively.

Although no special restriction is meant, when one pair of abstract data are to be processed in the shift operation unit 12, the information of the tag field to be conserved in a processed result is set to be the information of the tag field TFb which is latched in one input latch circuit ILAT2B. In other words, data which contains the information of the tag field to be conserved in the shift operation result is supplied to the input latch circuit ILAT2B.

In Fig. 5, symbol REARR denotes rearrange means for rearranging the information items of the data fields DFa and DFb respectively contained in one pair of input data to-be-processed, into a serial form.

Although no special restriction is intended, the rearrange means REARR operates in the case of the input data  $A_0 - A_{15}$  and  $B_0 - B_{15}$  being the data with the tags, to change the arrayal of the respective bits into the order of the data fields DFb and DFa and the tag fields TFb and TFa and to deliver the rearranged bits as outputs, as illustrated in Fig. 6. In a case where the input data items  $A_0 - A_{15}$  and  $B_0 - B_{15}$  are the data with no tag, the rearrange means does not rearrange the respective bits of the input data and delivers them as outputs in an order left intact.

Although a circuit arrangement for such rearrangement of the input bits is not especially restricted, it includes as shown in Fig. 6, a clocked inverter array CINVA1 which receives the bits  $B_{15} - B_{12}$  and which are coupled to the input terminals of output inverters  $INV_{31} - INV_{23}$ , a clocked inverter

array CINVA2 which receives the bits  $A_{15} - A_{12}$  and which are coupled to the input terminals of output inverters  $INV_{15} - INV_{12}$ , a clocked inverter array CINVA3 which receives the bits  $B_{15} - B_{12}$  and the bits  $A_{15} - A_{12}$  and which are coupled to the input terminals of output inverters  $INV_7 - INV_0$ , a clocked inverter array CINVA4 which receives the bits  $B_{11} - B_0$  and which are coupled to the input terminals of output inverters  $INV_{27} - INV_{16}$ , a clocked inverter array CINVA5 which receives the bits  $A_{11} - A_0$  and which are coupled to the input terminals of output inverters  $INV_{11} - INV_0$  and a clocked inverter array CINVA6 which receives the bits  $B_{11} - B_0$  and  $A_{11} - A_0$  and which are coupled to the input terminals of the output inverters  $INV_{31} - INV_8$ .

A control signal  $\phi_d$  is supplied to the selection terminals of the clocked inverter arrays CINVA1, CINVA2, CINVA4 and CINVA5, while a control signal  $\phi_t$  is supplied to the selection terminals of the clocked inverter arrays CINVA3 and CINVA6. Although not especially restricted, the control signals  $\phi_d$  and  $\phi_t$  are output from the microinstruction decoder 7 which decodes microinstructions read out of the microinstruction ROM 4 in succession. The series of microinstructions for each of various shift operations as stored in the microinstruction ROM 4 contain information for specifying the sort of the corresponding shift operation, and information for designating whether or not data to be processed are data with tags. In accordance with the above information items, in a case where the data to be processed are the data with the tags, the control signal  $\phi_t$  is brought to a select level such as high level, whereby the input bits have their arrayal changed and are then output as described above. In a case where the data to be processed are the data with no tag, the control signal  $\phi_d$  is brought to a select level such as high level, whereby the bits of the input data are not rearranged and are output in the order left intact.

In Fig. 5, symbol SFT denotes shift operation means for executing the shift operation for the 32-bit data delivered as the outputs from the rearrange means REARR, thereby to deliver shift data  $S_{15} - S_0$  of 16 bits as outputs. Although not especially restricted, the shift operation means SFT is constructed of an output selecting switch array SSWA of 32 columns which can select and output the 16-bit shift data  $S_{15} - S_0$  in one aspect among, at most, 32 output aspects that are obtained by shifting the 32 bits of the input data one by one in the rightward direction as viewed in the figure, a shift counter SCUNT which determines the number of shift bits of the input data within a range up to 32 bits, and a decoder SDEC which decodes the outputs of the shift counter SCUNT to form selection signals  $SEL_0 - SEL_{31}$  for selecting one predetermined column among the 32 columns of the output



selecting switch array SSWA, and which delivers the output shift data  $S_{15} - S_0$  conforming to the shift number instructed by the shift counter SCUNT.

In Fig. 5, symbol REPL denotes replace means capable of replacing the upper 4 bits  $S_{15} - S_{12}$  of the shift data  $S_{15} - S_0$  delivered as the outputs from the output selecting switch array SSWA, with the information items  $B_{15} - B_{12}$  contained in the tag field TFb of the input data  $B_{15} - B_0$ .

Although not especially restricted, the replace means REPL is constructed as a multiplexor by which the bits  $S_{15} - S_{12}$  or the bits  $B_{15} - B_{12}$  are selectively output as illustrated in Fig. 7. More specifically, it includes AND gates  $ANDd_{15} - ANDd_{12}$  which receive the bits  $S_{15} - S_{12}$  and the control signal  $\phi_d$  as 2 inputs, AND gates  $ANDt_{15} - ANDt_{12}$  which receive the bits  $B_{15} - B_{12}$  and the control signal  $\phi_t$  as 2 inputs, and OR gates  $OR_{15} - OR_{12}$  which receive the outputs of the AND gates  $ANDd_{15} - ANDd_{12}$  and AND gates  $ANDt_{15} - ANDt_{12}$  as 2 inputs. Thus, under the action of the control signal  $\phi_t$  which is controlled to the high level in response to the shift operation process of the data with the tags, the replace means REPL replaces the upper 4 bits  $S_{15} - S_{12}$  of the shift data  $S_{15} - S_0$  with the information items  $B_{15} - B_{12}$  contained in the tag field TFb of the input data  $B_{15} - B_0$  and then affords these information items to an output latch circuit OLAT2, whereas under the action of the control signal  $\phi_d$  which is controlled to the high level in response to the shift operation process of the data with no tag, the replace means affords the upper 4 bits  $S_{15} - S_{12}$  of the shift data  $S_{15} - S_0$  to the output latch circuit OLAT2 without any change.

Next, the operation of the shift operation unit 12 will be described.

First, the case where the input data items  $A_{15} - A_0$  and  $B_{15} - B_0$  are the data with the tags will be explained with reference to Fig. 8.

In the case where the input data items  $B_0 - B_{15}$  and  $A_0 - A_{15}$  to be subjected to the shift operation are the data with the tags, the control signal  $\phi_t$  which is output from the microinstruction decoder 7 is responsively set at the high level, and the control signal  $\phi_d$  at a low level. Thus, the rearrange means REARR changes the array of the respective bits into the order of the data field DFb, data field DFa, tag field TFb and tag field TFa so as to put the information items of the data fields DFa and DFb of both the input data into the serial form and then delivers the rearranged bits as outputs.

The shift operation means SFT receiving the outputs of the rearrange means REARR shifts the input data thereof a predetermined number of bits in accordance with a microprogram control, and delivers the shift data  $S_{15} - S_0$  of 16 bits as outputs. By way of example, when the shift number

is set at 6 bits as illustrated in Fig. 8, the 16 bits of the shift data  $S_{15} - S_0$  are brought to  $B_5, \dots B_0, A_{11}, \dots A_2$ .

Among the shift data items  $S_{15} - S_0$  ( $B_5, \dots B_0, A_{11}, \dots A_2$ ), the upper 4 bits  $S_{15} - S_{12}$  ( $B_5 - B_2$ ) are supplied to the replace means REPL. Since, however, the control signal  $\phi_t$  is set at the high level in conformity with the process of the data with the tags, the replace means REPL selects and outputs the information items  $B_{15} - B_{12}$  of the tag field TFb. Thus, the tag field TFe of operated result data to be afforded to the output latch circuit OLAT2 conserves the information items  $B_{15} - B_{12}$  contained in the tag field TFb of the input data  $S_{15} - S_0$ , and the data field DFe thereof contains serial shift operation results ( $B_1, B_0, A_{11}, \dots A_2$ ) concerning the information items of the data fields DFa and DFb in the input data  $A_{15} - A_0$  and  $B_{15} - B_0$ .

The case where the input data items  $A_{15} - A_0$  and  $B_{15} - B_0$  are the data with no tag will be explained with reference to Fig. 9.

In the case where the input data items  $B_0 - B_{15}$  and  $A_0 - A_{15}$  to be subjected to the shift operation are the data with no tag, the control signal  $\phi_t$  which is output from the microinstruction decoder 7 is responsively set at a low level, and the control signal  $\phi_d$  at the high level. Thus, the rearrange means REARR does not rearrange the respective bits of the input data  $B_{15} - B_0$  and  $A_{15} - A_0$  and delivers them as outputs in the order left intact.

On this occasion, the shift operation means SFT receiving the outputs of the rearrange means REARR shifts the input data thereof a predetermined number of bits in accordance with a microprogram control, and delivers the shift data  $S_{15} - S_0$  of 16 bits as outputs. By way of example, when the shift number is set at 6 bits as illustrated in Fig. 9, the 16 bits of the shift data  $S_{15} - S_0$  are brought to  $B_9, \dots B_0, A_{15}, \dots A_{10}$ .

Among the shift data items  $S_{15} - S_0$  ( $B_9, \dots B_0, A_{15}, \dots A_{10}$ ), the upper 4 bits  $S_{15} - S_{12}$  ( $B_9 - B_6$ ) are supplied to the replace means REPL. Since, however, the control signal  $\phi_d$  is set at the high level in conformity with the process of the data with no tag, the replace means REPL selects the outputs of the upper 4 bits  $S_{15} - S_{12}$  ( $B_9 - B_6$ ) among the shift data. Thus, the shift data items  $S_{15} - S_0$  ( $B_9, \dots B_0, A_{15}, \dots A_{10}$ ) are afforded to the output latch circuit OLAT2 without any change.

According to the embodiments described above, the following effects can be attained:

The arithmetic and logic operation unit 11 included in the coprocessor 1 masks the tag field TFa of input data  $A_0 - A_{15}$  on one side into logic "0" for a logical sum, exclusive "or" or addition operation or the like and into logic "1" for a logical product or subtraction operation or the like, whereupon it executes the arithmetical or logical opera-

tion for the pair of input data, and on that occasion, it prevents a carry from lower bits from being transmitted to the upper 4 bits  $D_{15} - D_{12}$  of an operated result. Thus, the result of the arithmetical or logical operation between the data fields DFa and DFb of the input data can be obtained in the operated result  $D_0 - D_{15}$  while the information items  $B_{15} - B_{12}$  of the tag field TFb in the input data  $B_0 - B_{15}$  is conserved without any change. Accordingly, in executing an arithmetical or logical operation for data with tags, the step of separating data fields and tag fields from the data to-be-processed and the step of affixing a predetermined tag field to an operated result for the information items of the separated data fields are dispensed with, to bring forth the effect that the arithmetical or logical operation of the data with the tags can be executed at high speed by the step.

In the shift operation of data with tags, the shift operation unit 12 included in the coprocessor 1 rearranges the information items of data fields DFa and DFb respectively contained in the input data  $A_{15} - A_0$  and  $B_{15} - B_0$  by the use of rearrange means REARR so as to make them interlacingly serial and then subjects the serial information items to the shift operation, and as regards the results  $S_{15} - S_0$  of the operation, it replaces the upper 4 bits  $S_{15} - S_{12}$  with the information items  $B_{15} - B_{12}$  of a tag field TFb in the input data  $B_{15} - B_0$ , whereby these information items of the tag field TFb can be conserved in the result of the shift operation. Accordingly, in executing a shift operation for data with tags, the step of separating data fields and tag fields from the data to-be-processed and the step of affixing a predetermined tag field to an operated result for the information items of the separated data fields are dispensed with, to bring forth the effect that the shift operation of the data with the tags can be executed at high speed by one step.

Besides, regarding data with no tag as made up of only data fields, the arithmetic and logic operation unit 11 can release the masking function for the input data and the carry inhibiting function for the predetermined bits in the case of the arithmetical or logical operation, whereby it can execute the arithmetical and logical operations, not only for the data with the tags, but also for the data with no tag.

Similarly, regarding data with no tag, the shift operation unit 12 does not select the function of rearranging the input data by the use of rearrange means REARR and the function of replacing shift data by the use of replace means REPL, whereby it can execute the shift operations, not only for the data with the tags, but also for the data with no tag.

Accordingly, the execution unit 3 executes the operations of the data with the tags at high speed

and is also permitted to execute the operations of the data with no tag, thereby to contribute to enhancement in the performance of the run of an AI language program as the whole system including the coprocessor 1.

Although, in the above, the invention made by the inventors has been concretely described in conjunction with embodiments, it is needless to say that the present invention is not restricted to the foregoing embodiments, but that it can be variously altered within a scope not departing from the purport thereof.

In the embodiments, as conservation means for conserving the information of the predetermined second information field of input data in processed result output data, mask means MASK has been adopted in an arithmetic and logic operation unit, and rearrange means REARR and replace means REPL have been adopted in a shift operation unit, but the present invention is not restricted to them. By way of example, in the arithmetic and logic operation unit, the conservation means can be constructed in such a way that the replace means as explained in the shift operation unit is disposed at a stage preceding an output latch circuit. Further, the operation unit of the present invention is not restricted to the shift operation unit and the arithmetic and logic operation unit explained in the embodiments, but it can be constructed as an adder, a multiplier, etc.

In the embodiments, there have been described cases where the tag field to be conserved is contained in the upper bits of data containing this tag field. However, for data of a structure in which the tag field is contained in the lower bits of the data, operation units can be constructed in such a way that the positions and circuit arrangements of the conservation means as explained in the embodiments are brought into correspondence with the lower bit side of the input data. Further, the conservation means can be constructed so as to be selectively applicable to both the structure having the tag field on the lower bit side and the structure having the tag field on the upper bit side.

The rearranging form in which data fields are serialized by the rearrange means as explained in the embodiments, is a mere example, and the respective bits of one pair of input data may well be rearranged so as to be alternately arrayed. By way of example, in a case where the input data items  $A_{15} - A_0$  and  $B_{15} - B_0$  are data with no tag, they are not, in effect, subjected to the rearrangement of serializing the data fields thereof, and the respective input bits are arrayed as  $B_{15}, A_{15}, \dots, A_0, B_0$ , which are output. In a case where the input data items are data with tags, they are, in effect, subjected to the rearrangement of serializing the data fields thereof, and the respective input bits are

arrayed as B<sub>15</sub>, B<sub>3</sub>, B<sub>14</sub>, B<sub>2</sub>, B<sub>13</sub>, B<sub>1</sub>, B<sub>12</sub>, B<sub>0</sub>, A<sub>15</sub>, A<sub>11</sub>, A<sub>14</sub>, A<sub>10</sub>, A<sub>13</sub>, A<sub>9</sub>, A<sub>12</sub>, A<sub>8</sub>, B<sub>11</sub>, A<sub>7</sub>, B<sub>10</sub>, A<sub>6</sub>, ... B<sub>5</sub>, A<sub>1</sub>, B<sub>4</sub>, A<sub>0</sub>, which are output. Such rearranging systems can be brought into various aspects in conformity with the ways of selecting switch columns in shift operation means.

The embodiments have been described concerning abstract data the first information field of which is the data field and the second information field of which is the tag field for indicating the attribute of the data field, but the relation between the first information field and the second information field is not restricted thereto and can be properly altered. Further, the bit formats of the first information field and the second information field are not restricted to those of the embodiments.

In addition, the embodiments have been described as the operation units applicable to both the data with the tags and the data with no tag, but it suffices for the present invention that data of a structure having, at least, the first information field and the second information field can be processed. Besides, controls for the operation units are not restricted to microprogram controls, but they can also be performed by a wired logic control system.

Although, in the above, the invention made by the inventors has been chiefly described in relation to the cases of applications to processors handling an AI language as form the background field of utilization, the present invention is not restricted thereto but is extensively applicable to data processors such as general-purpose processors and various peripheral controllers.

## Claims

1. A data processor comprising execution means for executing instructions, said execution means including operation means capable of processing data of a structure which has a first information field and a second information field, said operation means including conservation means for conserving information of the second information field of the data to-be-processed in output data of an operated result.

2. A data processor according to Claim 1, wherein said conservation means is mask means for masking the second information field contained in the input data on one side to-be-processed, into a predetermined logical value in accordance with a sort of a logical operation and then supplying the masked information field to logical operation means.

3. A data processor according to Claim 2, wherein said mask means has a control state in which the second information field of the input data

on one side is masked into logic "0" in accordance with any of logical sum, exclusive "or" and addition operations, and a control state in which the second information field of the input data on one side is masked into logic "1" in accordance with either of logical product and subtraction operations.

4. A data processor according to Claim 3, wherein said mask means includes a control state in which the masking is not performed for data of a structure made up of only a single information field.

5. A data processor according to Claim 1, wherein said conservation means includes: rearrange means for rearranging information items of the first information fields respectively contained in a pair of input data to-be-processed, into a serial form and then delivering the serial information items as outputs to shift operation means, and replace means capable of replacing part of data to be delivered as outputs from said shift operation means, with information of the second information field contained in the input data on one side.

6. A data processor according to Claim 5, wherein said rearrange means has a control state in which the rearrangement of information fields is not performed for a pair of input data of a structure made up of only a single information field, thereby to supply said shift operation means with the input data on both sides as left intact.

7. A data processor according to Claim 5, wherein said replace means has a control state in which, in response to a control state of said rearrange means, the replacing operation is not performed for input data of a structure made up of only a single information field.

8. A data processor comprising: controlled memory means for storing instructions to-be-executed, and execution means for executing the instructions successively read out of said controlled memory means,

said execution means including operation means capable of processing data of a structure which has a data field and a tag field expressive of either of an attribute and a type of the data field,

said operation means including logical operation means, and mask means for masking the tag field contained in the input data on one side to-be-processed, into a predetermined logical value in accordance with a sort of a logical operation and then supplying the masked tag field to said logical operation means.

9. A data processor according to Claim 8, wherein said operation means includes: rearrange means for rearranging information items of the data fields respectively contained in a pair of input data to-be-processed, into a serial form and then delivering the serial information items as outputs,

shift operation means for shifting the data supplied from said rearrange means, at will and then delivering the shifted data as outputs, and replace means capable of replacing part of data to be delivered as outputs from said shift operation means, with information of the tag field contained in the input data on one side.

10. A data processor according to Claim 9, wherein:

said mask means has a control state in which the tag field of the input data on one side is masked into logic "0" in accordance with any of logical sum, exclusive "or" and addition operations, a control state in which the tag field of the input data on one side is masked into logic "1" in accordance with either of logical product and subtraction operations, and a control state in which the masking is not performed for data of a structure made up of only the data field,

said rearrange means has a control state in which the rearrangement is not performed for the pair of input data of the structure made up of only the data field, thereby to supply said shift operation means with the input data on both sides as left intact,

said replace means has a control state in which, in response to the control state of said rearrange means, the replacing operation is not performed for the input data of the structure made up of only the data field, and

said controlled memory means contains control information for selectively designating the control states in accordance with the sort of the operation and the sort of the data to-be-processed.

11. In a data processor having execution means for executing instructions;

a data processor characterized in that said execution means comprises operation means for executing an operation between at least two data items each of which is of a structure having a first information field and a second information field, and conservation means for conserving information of the second information field of at least one of the data items to-be-operated, in output data of an operated result.

12. In a data processor having controlled memory means for storing instructions to-be-executed, and execution means for executing the instructions successively read out of the controlled memory means;

a data processor characterized in that said execution means includes operation means for executing an operation between at least two data items each of which is of a structure having a data field and a tag field expressive of either of an attribute and a type of the data field,

wherein said operation means includes logical operation means, and mask means for masking the tag field contained in one of the data items to-be-

operated, into a predetermined logical value in accordance with a sort of the logical operation and then supplying the masked tag field to said logical operation means.

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FIG. 1

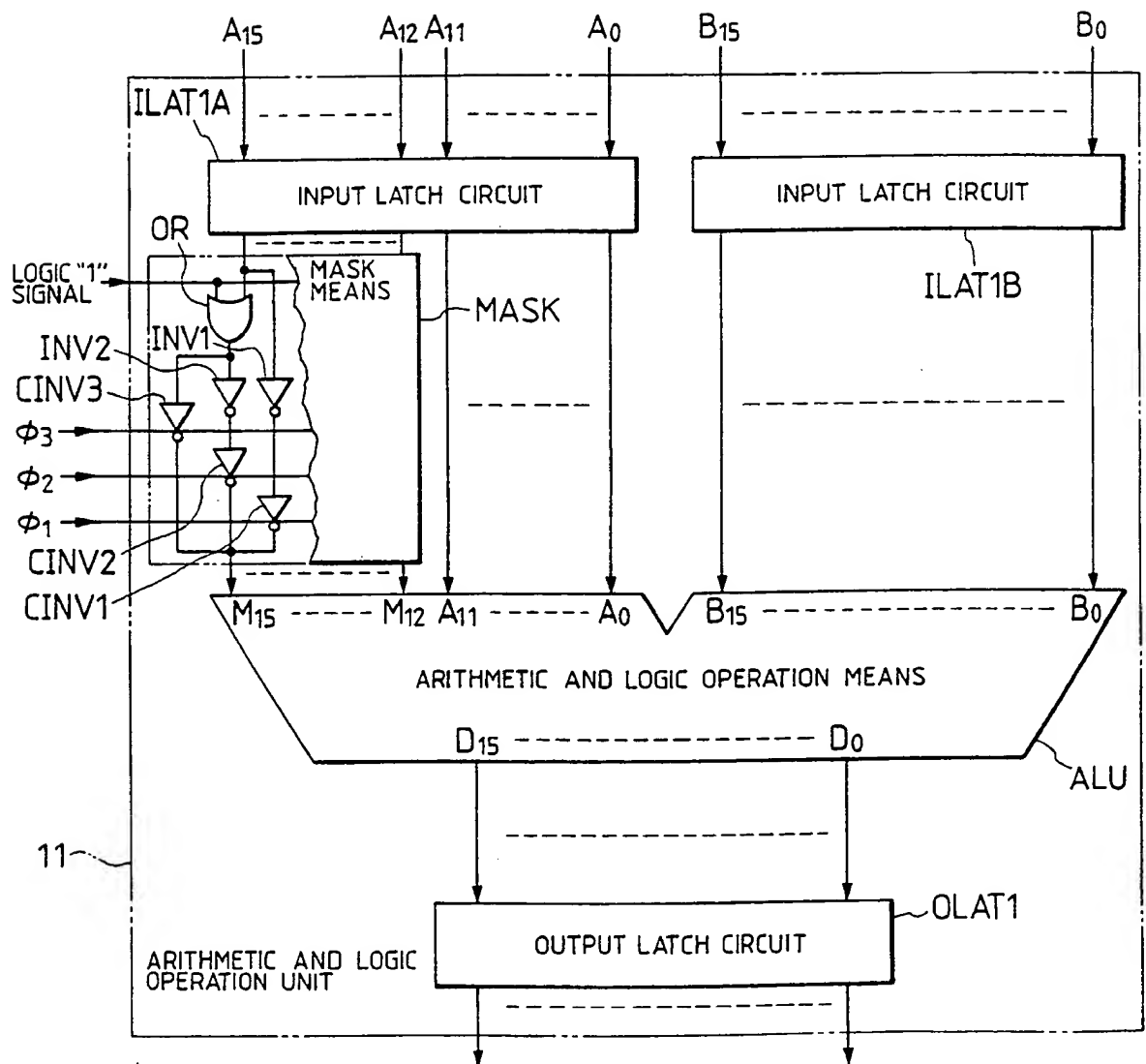


FIG. 2

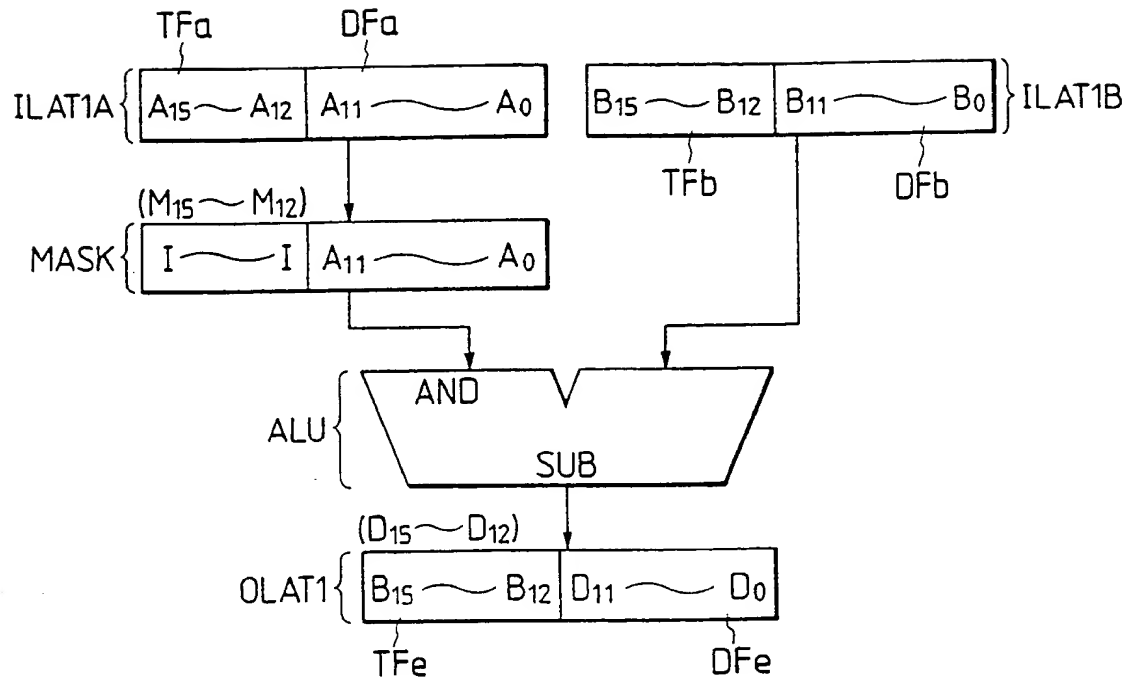


FIG. 3

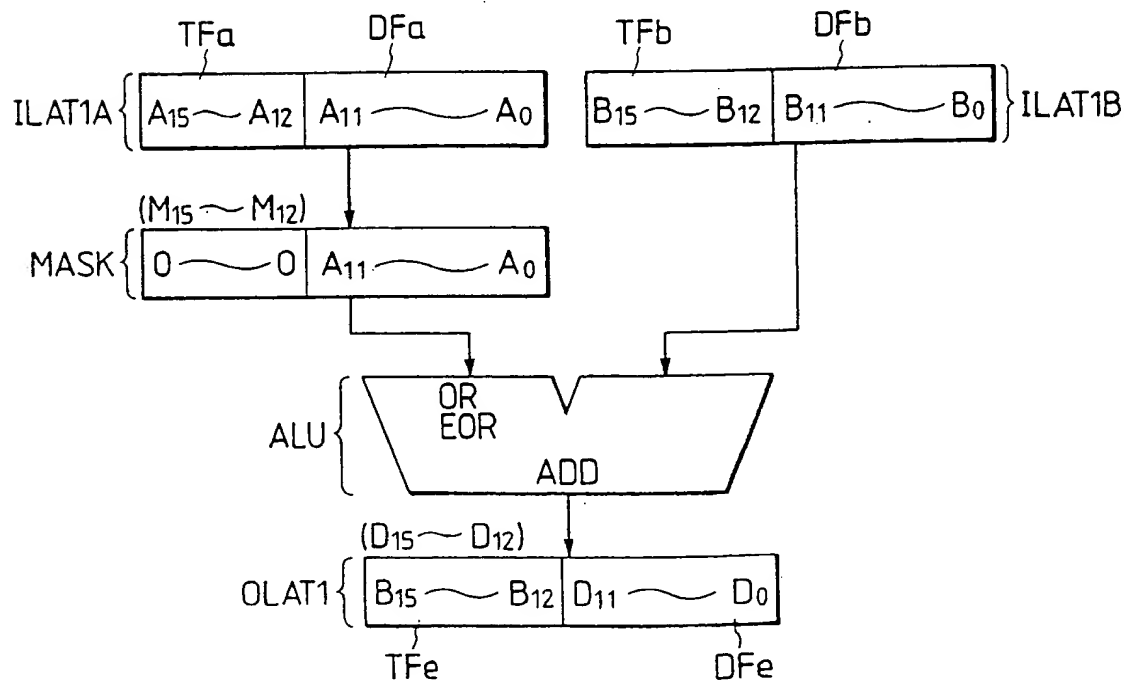


FIG. 4

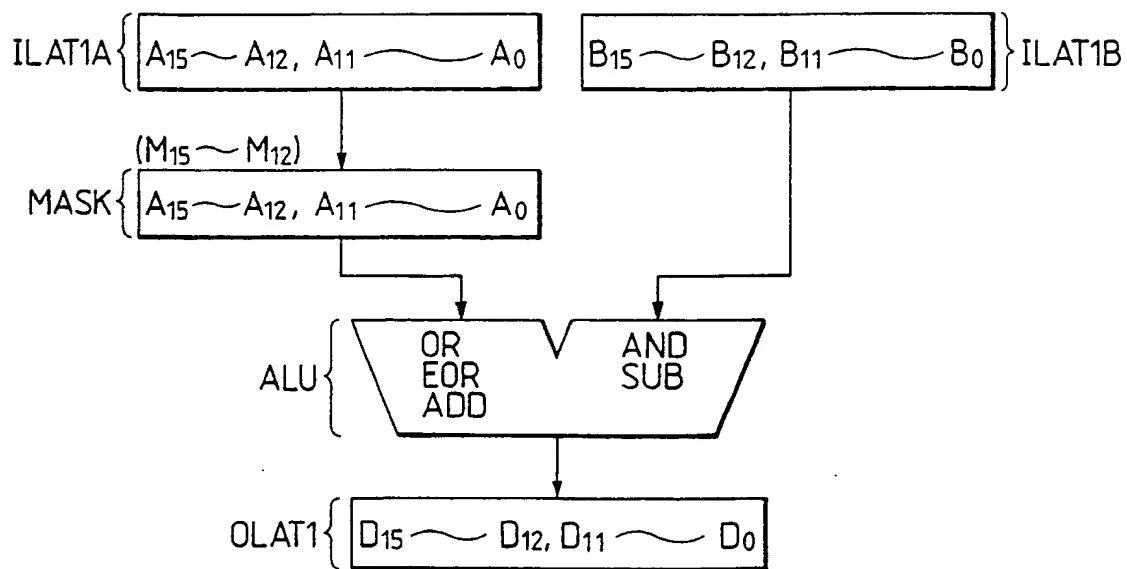
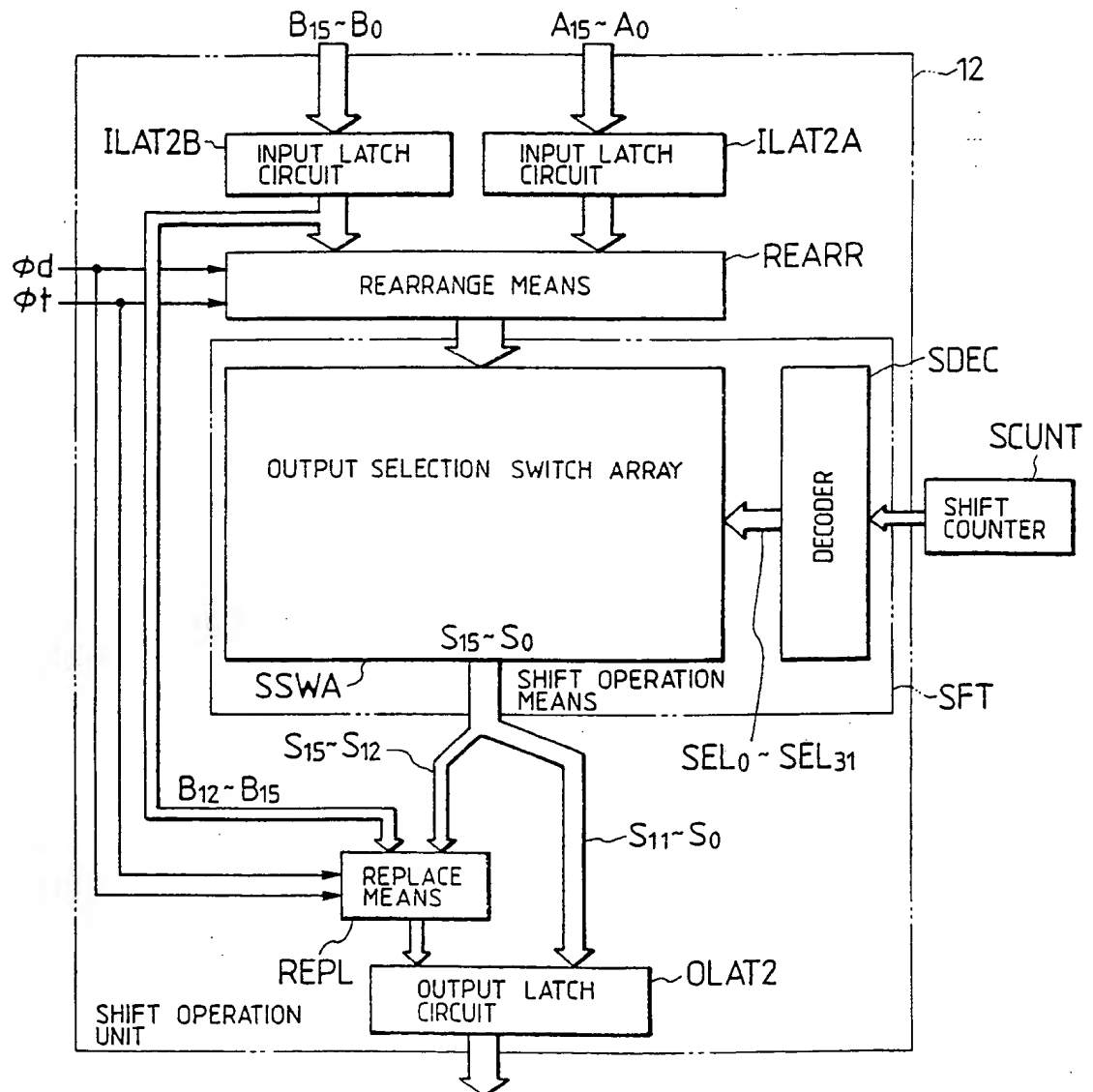


FIG. 5





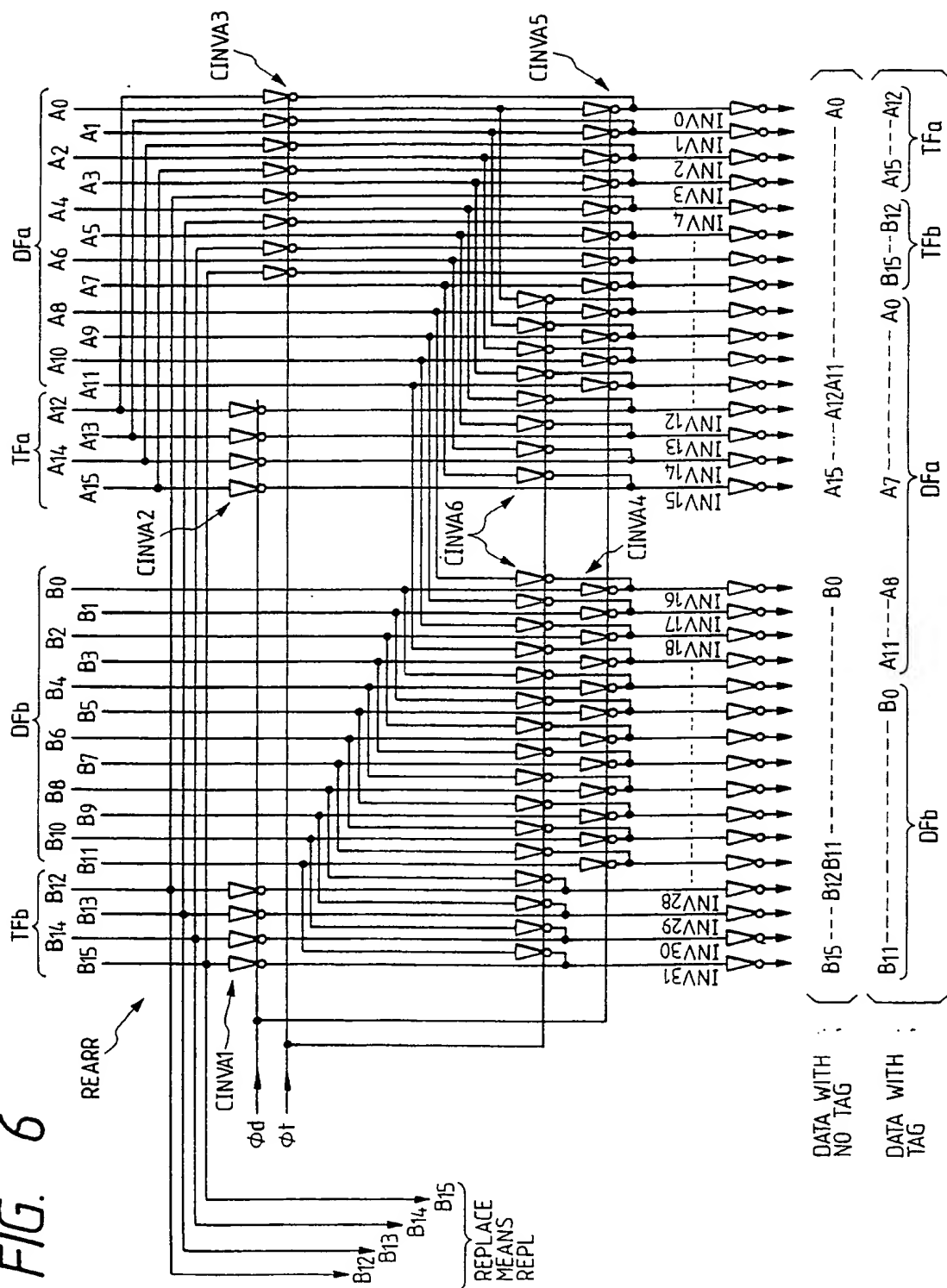


FIG. 7

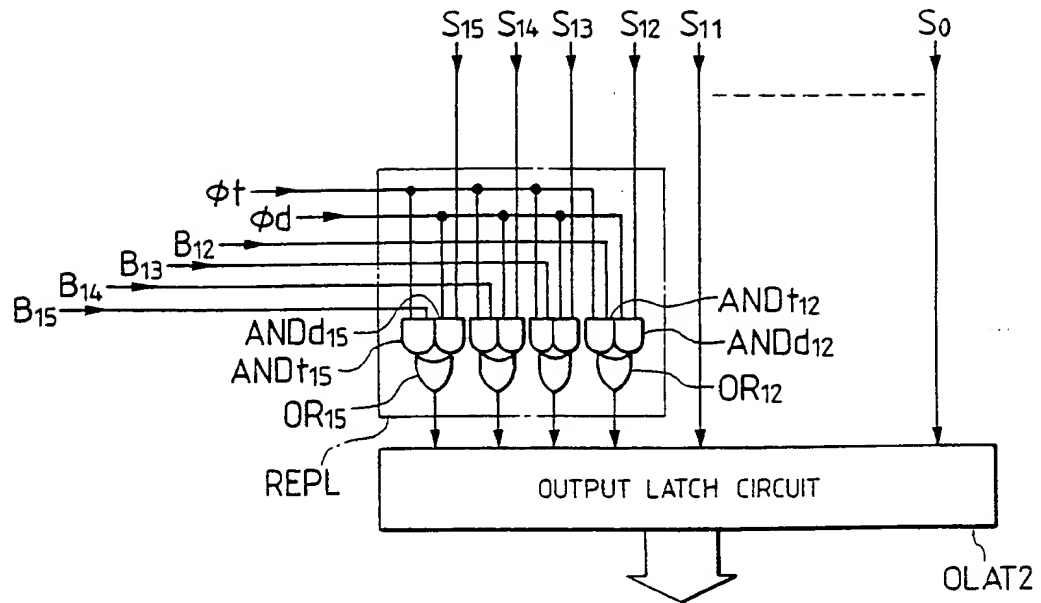


FIG. 8

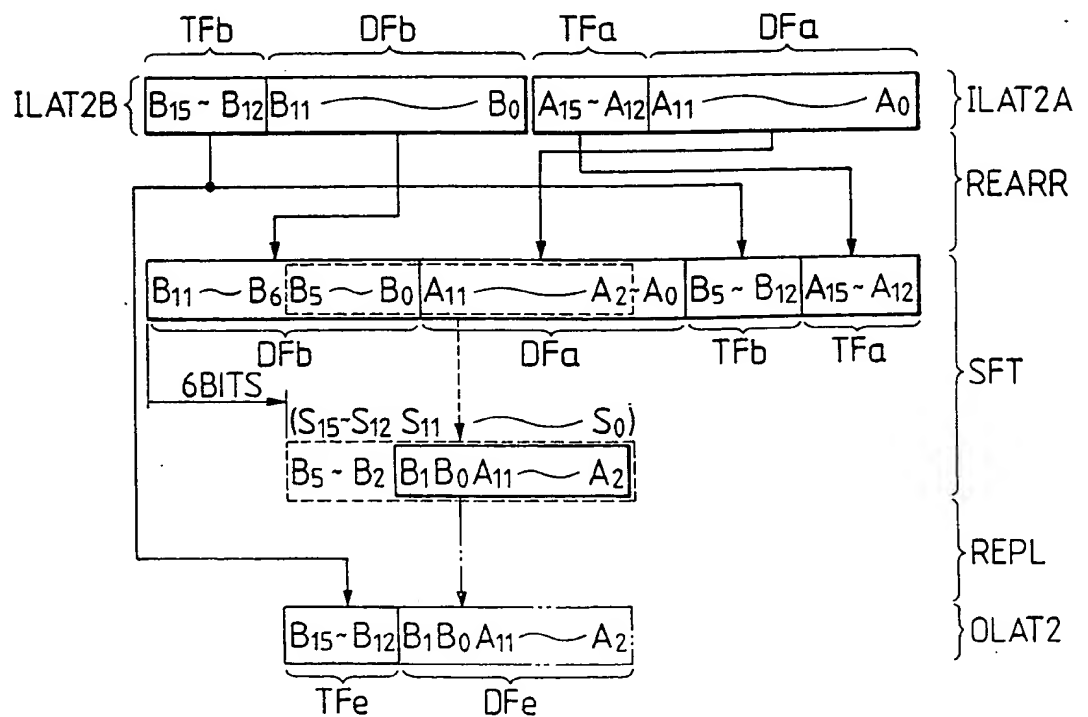


FIG. 9

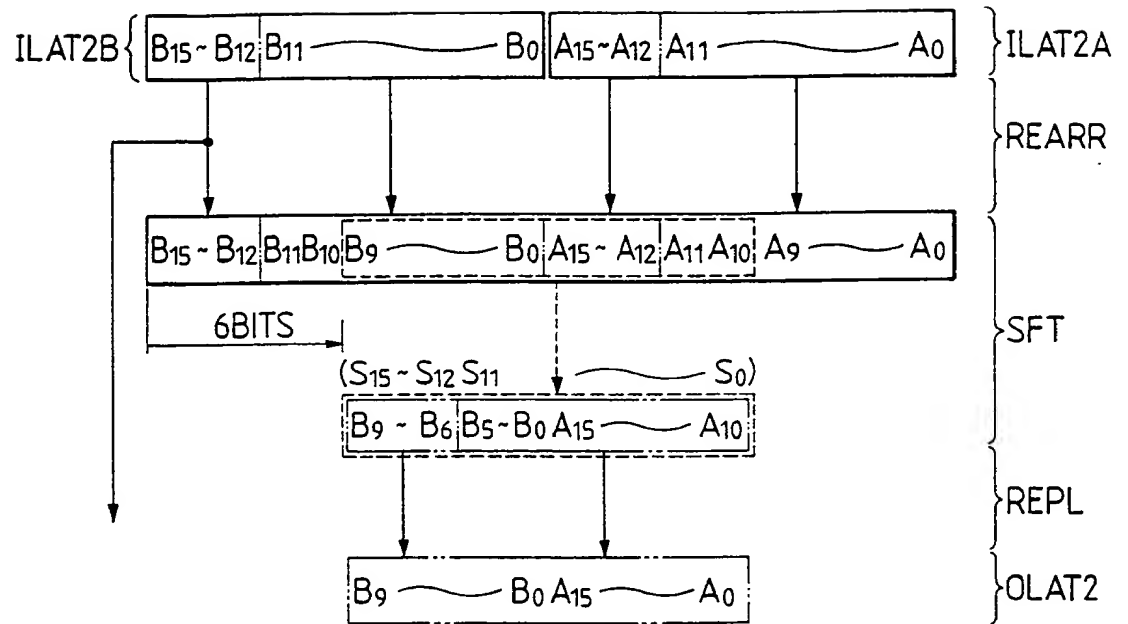


FIG. 10

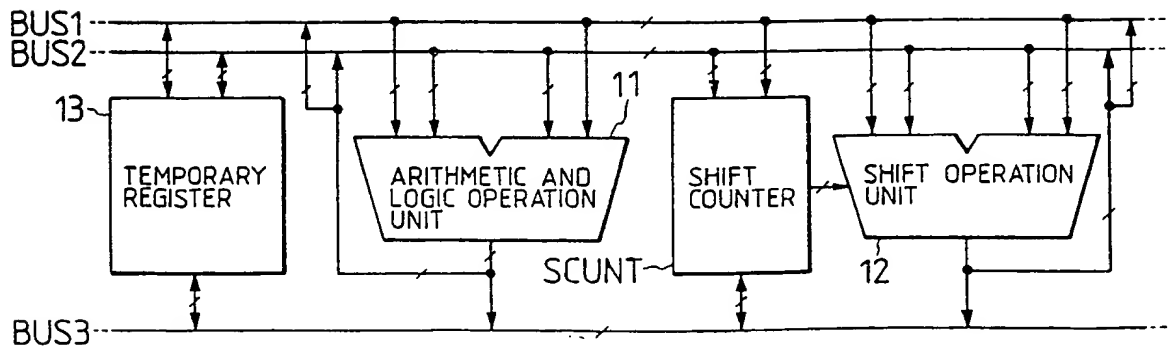
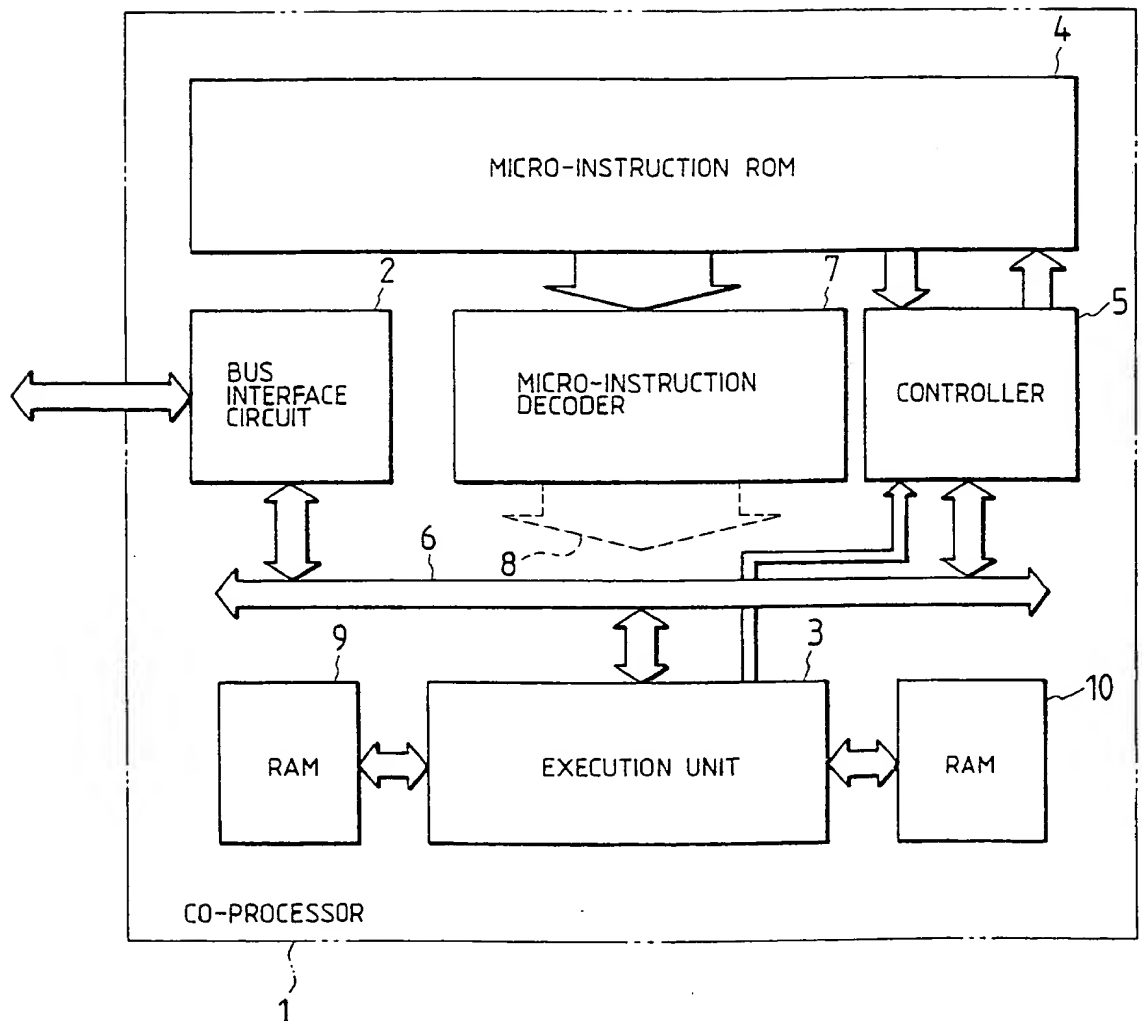


FIG. 11





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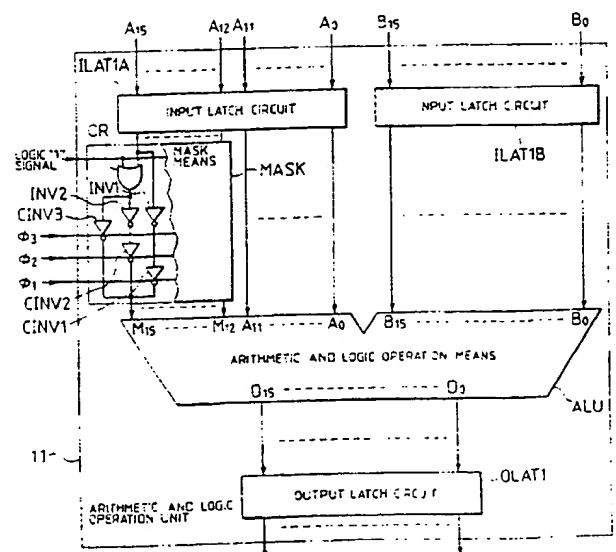
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Data processor.

A processor capable of processing data of a structure having a first information field and a second information field includes conservation means for conserving the information of the predetermined second information field of the input data. Thus, in executing a logical operation or a shift operation, the step of separating the first data fields and second data fields from the data to-be-processed and the step of afficing the predetermined second field to the result of the operation are dispensed with, thereby to achieve a high-speed operation for the data of the structure having the first information field and the second information field.

FIG. 1



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Patent Office

## EUROPEAN SEARCH REPORT

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DOCUMENTS CONSIDERED TO BE RELEVANT			
Category	Citation of document with indication, where appropriate, of relevant passages	Relevant to claim	CLASSIFICATION OF THE APPLICATION (Int. Cl.5)
A	EP-A-0 207 519 (HITACHI LTD) * column 7, line 51 - column 9, line 19 * * column 10, line 10 - column 11, line 44: figures * - - -		G 06 F 9:30 G 06 F 7 48
A	COMPUTER, vol. 20, no. 1, January 1987, LONG BEACH US pages 43 - 52; D.A.MOON: "Symbolics Architecture" * page 47, middle column, lines 59 - 62 * * page 48, middle column, lines 14 - 17 @ page 51, middle column, lines 3 - 4 * - - -		
P,A	EP-A-0 250 130 (ADVANCED MICRO DEVICES INC) * abstract: figure * * column 5, lines 8 - 54 * - - -		
A	EP-A-0 018 298 (ANVAR) * abstract * - - -		
A	FR-A-2 550 362 (CAZOR) * abstract: figure 3 * * page 16, line 30 - page 17, line 9 * - - - - -		
The present search report has been drawn up for all claims			TECHNICAL FIELDS SEARCHED (Int. Cl.5)
			G 06 F
Place of search		Date of completion of search	Examiner
The Hague		04 January 91	QUESSON C.J.
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